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respective date stamp files (the machine A, which has already been disconnected, does not receive the interrupt signal caused by the disconnection of B and C).

In figure 3G, the machines A and B again connect onto the bus, with A as reference machine, and thus pass on to the sections α3 and β4, whereas the machine C remains disconnected and, since it does not receive an interrupt signal, continues to fill out the section γ5.

Subsequently, figure 3H, the three machines are all again connected to the bus (sections α4, β5 and γ6), with B as reference machine.

When the latter (B) is disconnected (figure 3I, sections α5, β6 and γ7), A becomes the reference machine. A short time later (figure 3~~4~~J, sections α6, β7 and γ8), B reconnects and again becomes reference machine.

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In figure 3K, the machine C disconnects (sections α7, β8 and γ9) and A again becomes reference machine.

Finally (figure 3L, sections α8 and β9), the machines A and B are also disconnected, after which the history of the system becomes irrelevant.

Thanks to the second temporal reference (TMR), a kind of re-synchronization has thus been performed, which is however very rudimentary because the rows of the date stamp files are filled out at a slow rate. This bringing into correspondence of the sections α, β and γ allows the non-monotonicity of the first temporal reference (NT) to be overcome and allows it to be used to carry out the actual re-synchronization by intervals.

In this example, the case where only one bus is present has been considered.

Coming back to the problem of re-synchronizing the data of the machine B with respect to the time of the machine A, the 'common temporal ranges' must first be determined, in other words the periods during which these two machines were connected to the same network via the bus BUS, which is that of the only network considered in this example. There are four of these ranges: P1, which corresponds to the sections α3 and β4; P2,